

COLUMNS vs. ROWS

Influence of the Reduction Order in Multiplier Verification using Computer Algebra

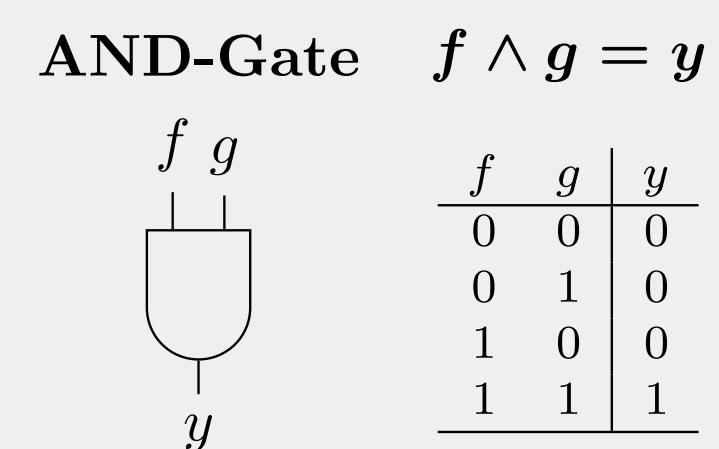
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Computer Algebra

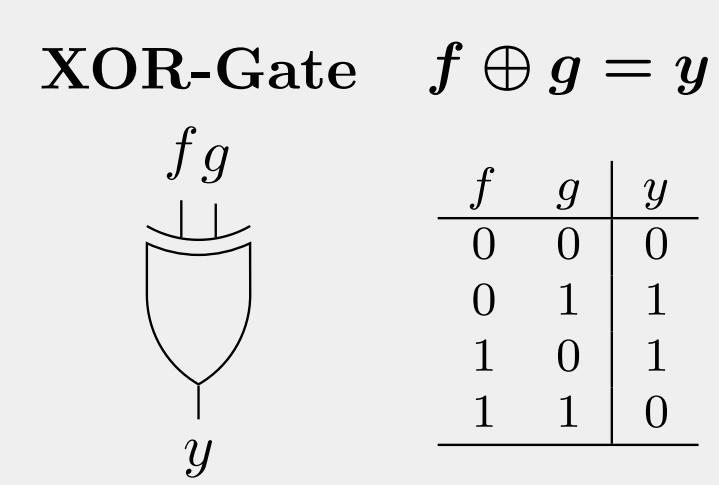
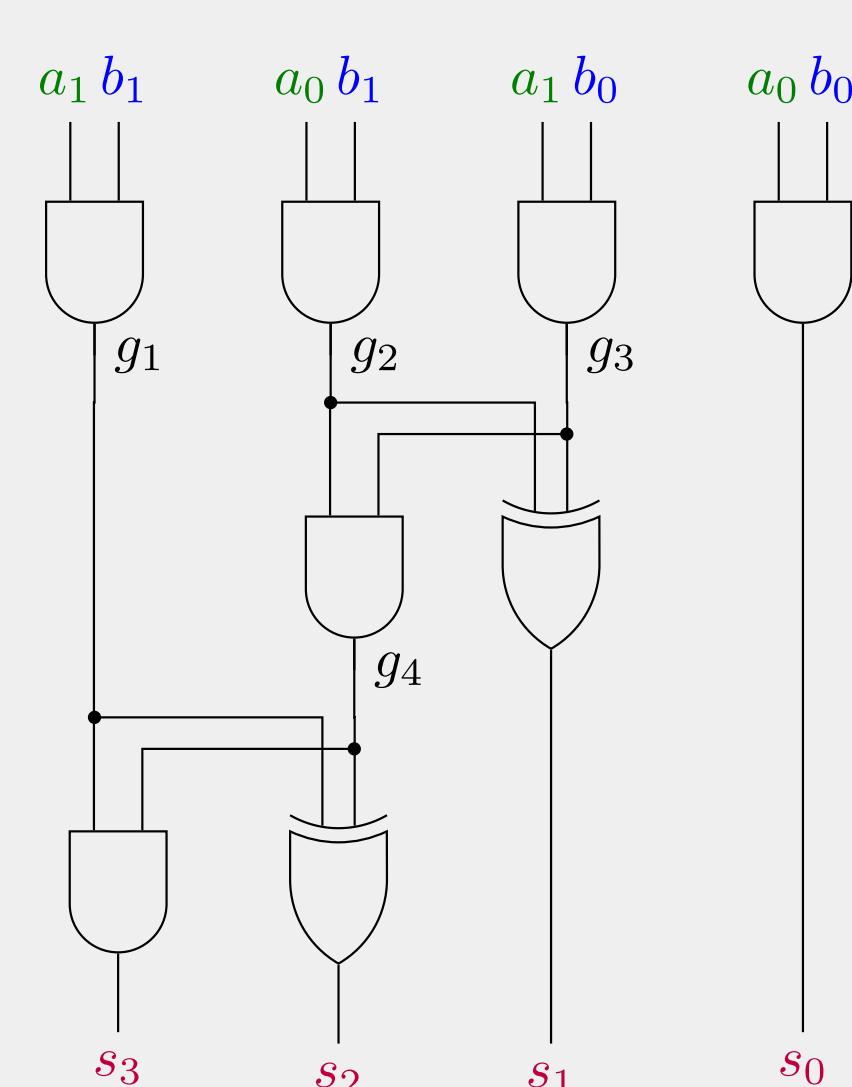
Bugs in hardware are expensive!

Challenge: Multiplication

System



$$y = fg$$



$$y = f + g - 2fg$$

Model

$$\begin{aligned} s_3 &= g_1 \wedge g_4 & -s_3 + g_1 g_4, \\ s_2 &= g_1 \oplus g_4 & -s_2 - 2g_4 g_1 + g_4 + g_1, \\ g_4 &= g_2 \wedge g_3 & -g_4 + g_2 g_3, \\ s_1 &= g_2 \oplus g_3 & -s_1 - 2g_2 g_3 + g_2 + g_3, \\ g_1 &= a_1 \wedge b_1 & -g_1 + a_1 b_1, \\ g_2 &= a_0 \wedge b_1 & -g_2 + a_0 b_1, \\ g_3 &= a_1 \wedge b_0 & -g_3 + a_1 b_0, \\ s_0 &= a_0 \wedge b_0 & -s_0 + a_0 b_0, \\ a_1 &\in \mathbb{B} & -a_1^2 + a_1, \\ a_0 &\in \mathbb{B} & -a_0^2 + a_0, \\ b_1 &\in \mathbb{B} & -b_1^2 + b_1, \\ b_0 &\in \mathbb{B} & -b_0^2 + b_0 \end{aligned}$$

Specification

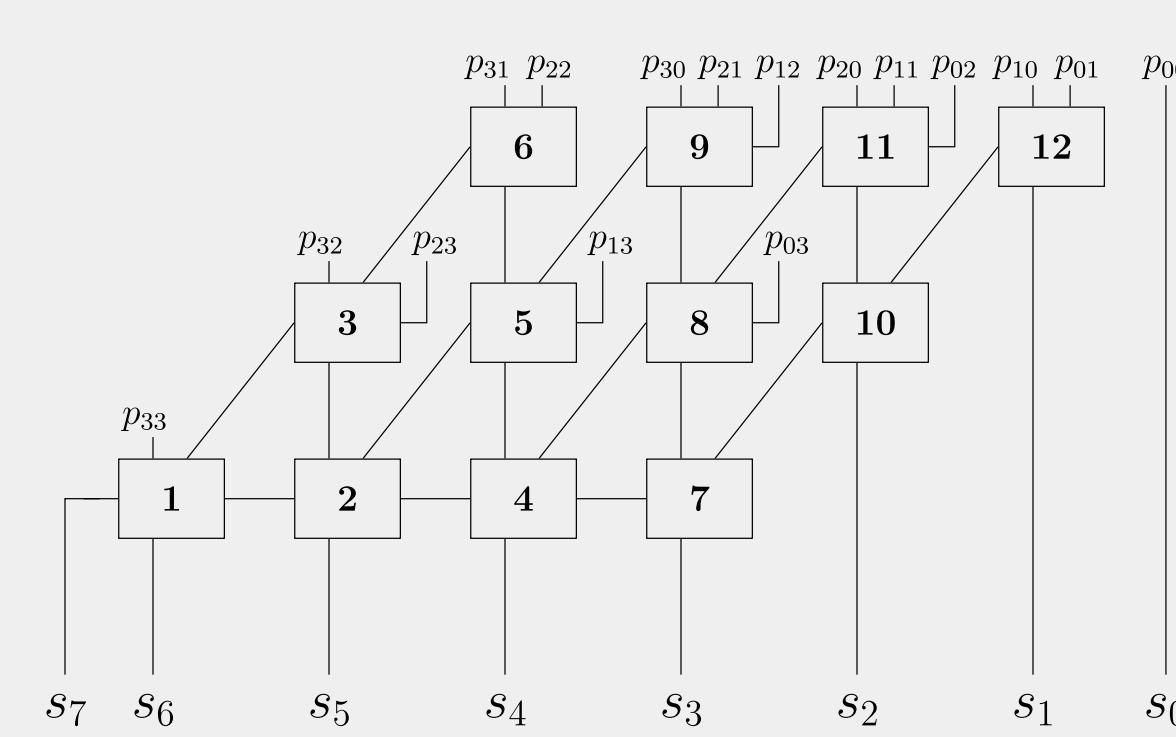
$$(2a_1 + a_0) * (2b_1 + b_0) = 8s_3 + 4s_2 + 2s_1 + s_0$$

Reduction

$$\begin{aligned} &8s_3 + 4s_2 + 2s_1 + s_0 - 4a_1 b_1 - 2a_1 b_0 - 2a_0 b_1 - a_0 b_0 \\ &4s_2 + 8g_4 g_1 + 2s_1 + s_0 - 4a_1 b_1 - 2a_1 b_0 - 2a_0 b_1 - a_0 b_0 \\ &4g_4 + 2s_1 + 4g_1 + s_0 - 4a_1 b_1 - 2a_1 b_0 - 2a_0 b_1 - a_0 b_0 \\ &2s_1 + 4g_2 g_3 + 4g_1 + s_0 - 4a_1 b_1 - 2a_1 b_0 - 2a_0 b_1 - a_0 b_0 \\ &\vdots \\ &0 \end{aligned}$$

Preprocessing eliminates internal nodes of full- and half-adders.
Specification is reduced by the rewritten gate polynomials.

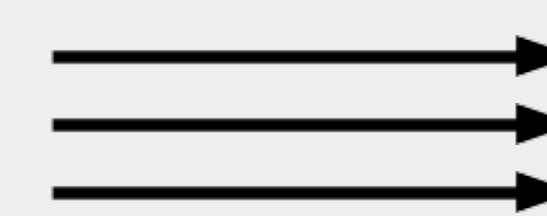
Polynomial reduction orderings:



Column-wise Order



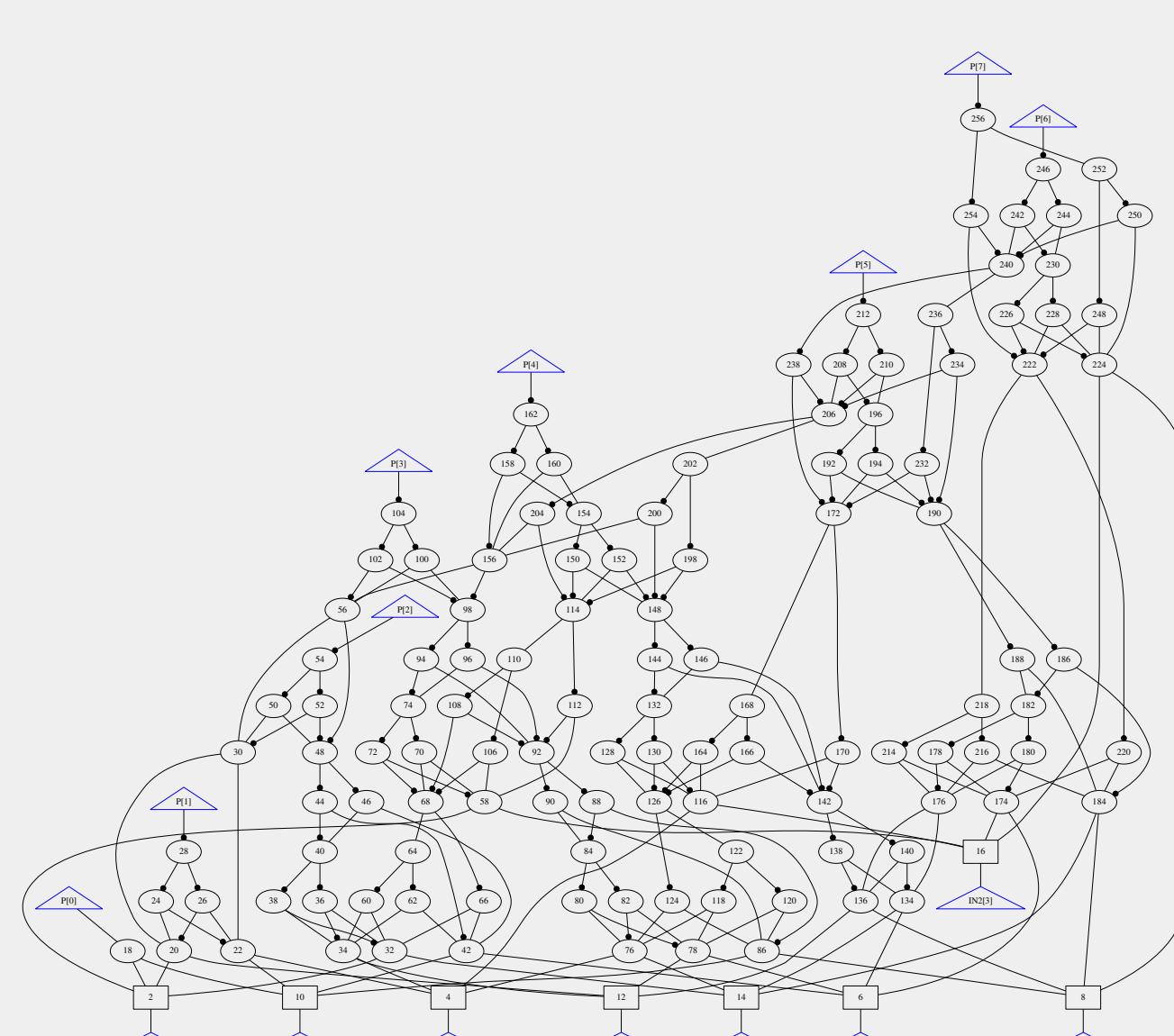
Row-wise Order



Arbitrary reverse topological



Simple Multiplier

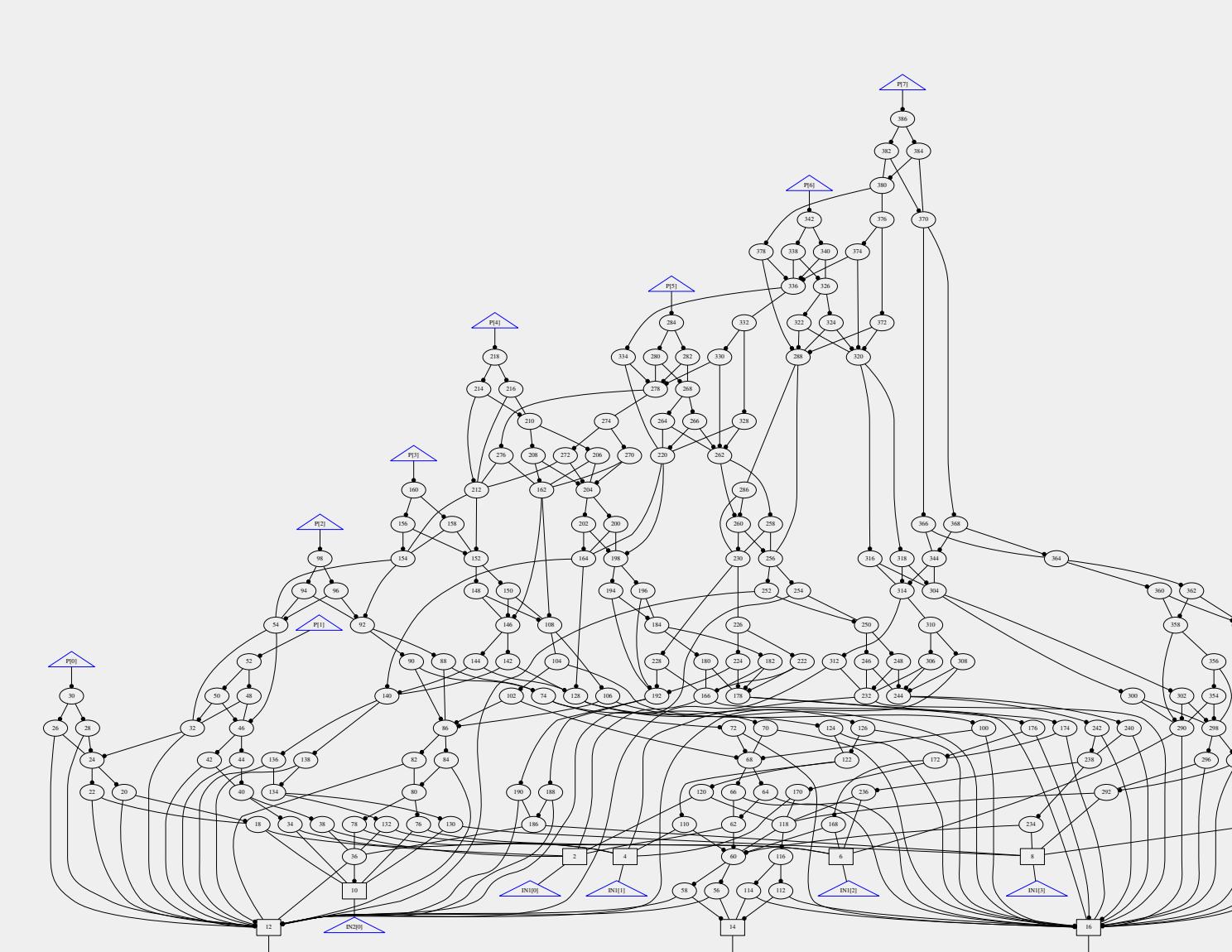


Final stage adder:
Ripple-Carry adder

PP accumulation:
Array

PP generation:
AND Gates

Complex Multiplier



Final stage adder:
Ripple-Carry adder

PP accumulation:
Wallace-Tree

PP generation:
Booth Encoding

